

* Labels. let L be a countable infinite set of labels. we write $l, l_1, l_2, \dots \in L$.

We extend notion of context with labels.

$$\Gamma ::= \emptyset \mid \Gamma, x:A \mid \Gamma, l: \text{Qubit}.$$

We write $\Sigma := [l_1: \text{Qubit}, \dots, l_n: \text{Qubit}]$.

We call Σ Label context.

* Labels corresponds to wires, or the addresses of the qubits.

* Values. we extend the notion of value.

$$V ::= \dots \mid l \mid \text{lift } M \mid \mu \mid \text{Init} \mid \text{Meas.} \mid \text{force } V.$$

* Label vs variable

A variable can be substituted by another term, but we can't substitute a label with another term.

~~we can only rename a label to another label~~

* Evaluation

Rather than evaluating only a term like STLC, we evaluate a "configuration".

A configuration is a triple (φ, L, M) .

where L is a finite list of labels $[l_1, \dots, l_n]$,

and $\varphi \in \underbrace{\text{Qubit} \otimes \dots \otimes \text{Qubit}}_{|L|}$

and M is a term.

We write φ_{l_i} for the i -th qubit in φ .

$$\text{e.g. } \varphi = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$$

φ_{l_1} refers to the first qubit of φ .

φ_{l_2} refers to the second qubit.

$$\textcircled{2} \varphi = \beta_{00} \otimes |0\rangle.$$

$$\varphi_{l_3} = |0\rangle.$$

* in (φ, L, M) , M may contain labels in L .

~~and~~

* Evaluation rules

These rules update the underlying quantum state.

$$(\varphi, L, \text{force } U)(l_i, \dots, l_j)$$

$$\frac{\varphi' = U_{l_i, \dots, l_j}(\varphi)}{(\varphi, L, \text{force } U)(l_i, \dots, l_j)} \rightarrow (\varphi', L, (l_i, \dots, l_j))$$

$$L' = L, l_{n+1}$$

$$(\varphi, L, \text{force Init}()) \rightarrow (\varphi \otimes |0\rangle, L', l_{n+1})$$

$$\begin{aligned} \varphi &= a\psi + b\gamma & \psi' &= \psi/|0\rangle \\ \psi_{l_i} &= |0\rangle & \gamma_{l_i} &= |1\rangle & \gamma' &= \gamma/|1\rangle \end{aligned}$$

$$(\varphi, L, \text{force meas } l_i) \xrightarrow{|a|^2} (\psi', L/l_i, \text{False})$$

...

$$(\varphi, L, \text{force meas } l_i) \xrightarrow{|b|^2} (\gamma', L/l_i, \text{True})$$

$$(\varphi, L, (\text{let } x.M) V) \rightarrow (\varphi, L, [V/x]M)$$

$$(\varphi, L, M) \rightarrow (\varphi')$$

$$(\varphi, L, MN) \rightarrow (\varphi', L', M')$$

$$(\varphi, L, \text{force } (l_i \# M)) \rightarrow (\varphi, L, M)$$

$$\begin{aligned} (\varphi, L, \text{let } (x, y) = (V_1, V_2) \text{ in } M) \\ \rightarrow (\varphi, L, [V_1/x, V_2/y]M) \end{aligned}$$