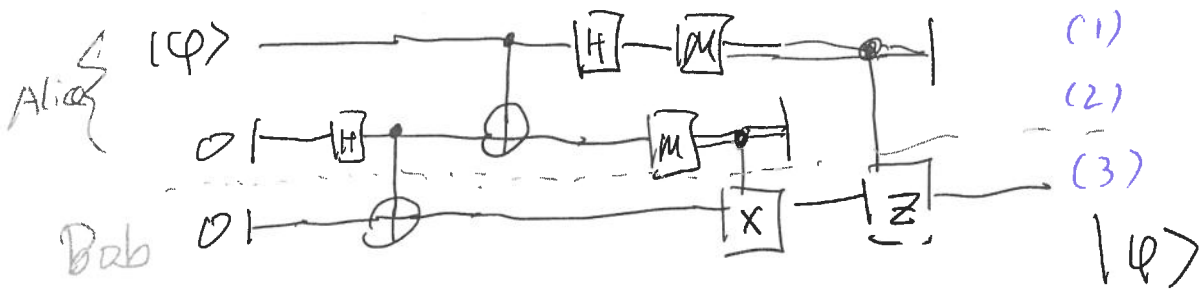


* Teleportation Circuit.



Teleportation protocol:

How would we verify this claim?

$$|\varphi\rangle|00\rangle \xrightarrow{B} \frac{1}{\sqrt{2}} (|\varphi\rangle|00\rangle + |\varphi\rangle|11\rangle)$$

$$\text{Let } |\varphi\rangle = a|0\rangle + b|1\rangle = \frac{1}{\sqrt{2}} (a|000\rangle + b|100\rangle + a|011\rangle + b|111\rangle)$$

$$\xrightarrow{CNOT_{1 \rightarrow 2}} \frac{1}{\sqrt{2}} (a|000\rangle + b|1110\rangle + a|011\rangle + b|1101\rangle)$$

$$\xrightarrow{H_1} \frac{1}{\sqrt{2}} (a|+\rangle|00\rangle + b|-\rangle|110\rangle + a|+\rangle|111\rangle + b|-\rangle|101\rangle)$$

$$= \frac{1}{2} (a|000\rangle + a|100\rangle + b|010\rangle - b|110\rangle + a|011\rangle + a|111\rangle + b|001\rangle - b|101\rangle)$$

$$= \frac{1}{2} \left(\begin{aligned} &|00\rangle (a|0\rangle + b|1\rangle) + \\ &|01\rangle (a|1\rangle + b|0\rangle) + \\ &|10\rangle (a|0\rangle - b|1\rangle) + \\ &|11\rangle (a|1\rangle - b|0\rangle) \end{aligned} \right)$$

So. $M_1 = M_2 = 0$. then we are done.

if $M_1 = 0$
 $M_2 = 1$, then we apply X gate
to correct.

if $M_1 = 1$
 $M_2 = 0$, then we apply Z gate
to correct.

otherwise, we apply  gate.

All these happen with equal probability.

* Principle of ~~defer~~ deferred measurement.

Measurements done in the middle of the circuit can be deferred to the end.

In particular, we have the following identity.



How to show this?

Suppose we have an input state $\varphi = a|00\rangle + b|01\rangle + c|10\rangle + d|11\rangle$.

LHS: (left hand side)

$$\begin{array}{l}
 \varphi \\
 \swarrow \searrow \\
 |a|^2 + |b|^2 \quad |c|^2 + |d|^2 \\
 \downarrow \quad \downarrow \\
 a|0\rangle + b|1\rangle \quad c|0\rangle + d|1\rangle
 \end{array}$$

$$\begin{array}{l}
 \downarrow U \\
 cU|0\rangle + dU|1\rangle
 \end{array}$$

so LHS = RHS.

RHS: (right hand side)

$$\begin{array}{l}
 \varphi \\
 \downarrow CU \\
 a|00\rangle + b|01\rangle + |1\rangle cU|0\rangle + |1\rangle dU|1\rangle \\
 \swarrow \searrow \quad \downarrow \\
 |a|^2 + |b|^2 \quad |c|^2 + |d|^2 \\
 \downarrow \quad \downarrow \\
 a|0\rangle + b|1\rangle \quad cU|0\rangle + dU|1\rangle
 \end{array}$$

* Construction of Toffoli gate.

Recall that Clifford gate set is

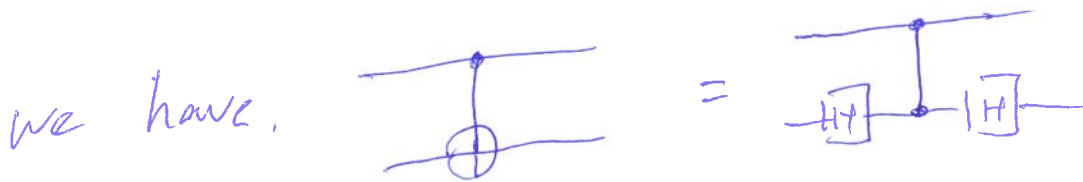
$$\{S, H, \text{CNOT}\}.$$

And $T|0\rangle = |0\rangle$
 $T|1\rangle = e^{i\frac{\pi}{4}}|1\rangle.$

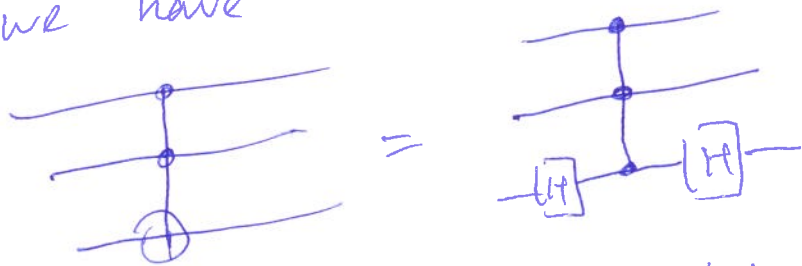
or $T|a\rangle = e^{i\frac{\pi}{4}a}|a\rangle, \text{ where } a \in \{0,1\}.$

So T moves the "state" a into the "phase" $e^{i\frac{\pi}{4}a}.$

$$\text{CNOT}|a,b\rangle = |a, a \oplus b\rangle, \text{ where } a, b \in \{0,1\}.$$



Similarly, we have



is CCZ gate (control-controlled-Z).

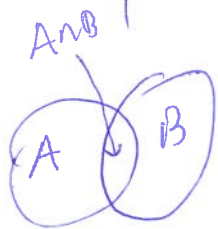
$$\text{CCZ}|a,b,c\rangle = e^{i\pi(a \cdot b \cdot c)}|a,b,c\rangle.$$

Note that $e^{i\pi} = -1.$

So to construct Toffoli gate using Clifford + T, we can just construct CCZ using Clifford + T.

* Inclusion-exclusion principle. (For set).

$$|A \cap B| = |A| + |B| - |A \cup B|$$



For binary number, we have:

$$2ab = a + b - a \oplus b$$

$$\text{So } 4abc = 2(2ab)c = 2(a + b - a \oplus b) \cdot c$$

$$= 2ac + 2bc - 2(a \oplus b)c$$

$$= a + c - a \oplus c + b + c - b \oplus c - (a \oplus b + c - a \oplus b \oplus c)$$

$$= a + b + c - a \oplus c - b \oplus c - a \oplus b + a \oplus b \oplus c$$

$$\text{So } e^{i\pi(abc)} = e^{i\frac{\pi}{4}(4abc)} = e^{i\frac{\pi}{4}(a+b+c - a \oplus c - b \oplus c - a \oplus b + a \oplus b \oplus c)}$$

$$= e^{i\frac{\pi}{4}a} \cdot e^{i\frac{\pi}{4}b} \cdot e^{i\frac{\pi}{4}c} \cdot e^{-i\frac{\pi}{4}(a \oplus c)} \cdot e^{-i\frac{\pi}{4}(b \oplus c)} \cdot e^{-i\frac{\pi}{4}(a \oplus b)} \cdot e^{i\frac{\pi}{4}(a \oplus b \oplus c)}$$