

Designing Quantum Programming Languages with Types

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Why quantum programming languages?

- ▶ Researchers have shown quantum algorithms can offer substantial speed-up for certain computing tasks.
- ▶ Advances in quantum hardware from companies like IBM and Google.
- ▶ Quantum algorithms are usually expressed using quantum circuits.
- ▶ Quantum algorithms are commonly expressed at a high level.
- ▶ Debugging quantum algorithms can be expensive.

My research interest

Build tools to facilitate programming quantum computers.

- ▶ How to design a high-level programming language for quantum circuits?
- ▶ How to verify quantum programs?
- ▶ How to run a high-level programming language on actual quantum computer?
- ▶ What algorithms to run on current quantum computer?

Why types?

- ▶ Lightweight specifications of programs.
- ▶ Allow compiler to enforce invariants via type checking.
- ▶ A well-typed program satisfies certain properties.

Background on types: an idealized programming language

- ▶ Programs $M, N := x \mid \lambda x.M \mid MN$.
- ▶ Types $A, B := C \mid A \rightarrow B$.
- ▶ Typing environment $\Gamma = x_1 : A_1, \dots, x_n : A_n$.
- ▶ Typing judgment $\Gamma \vdash M : A$.
- ▶ Typing rules

$$\frac{(x : A) \in \Gamma}{\Gamma \vdash x : A} \quad \frac{\Gamma, x : A \vdash M : B}{\Gamma \vdash \lambda x.M : A \rightarrow B} \quad \frac{\Gamma \vdash M : A \rightarrow B \quad \Gamma \vdash N : A}{\Gamma \vdash MN : B}$$

Type safety

- ▶ A *type checker* checks $\Gamma \vdash M : A$.
- ▶ An *evaluator* performs evaluation $M \Downarrow V$.
- ▶ *Type safety*
If $\Gamma \vdash M : A$ and $M \Downarrow V$, then $\Gamma \vdash V : A$.

Fancy types

- ▶ Linear types: $A \multimap B$.
- ▶ Dependent types: $(n : \text{Nat}) \rightarrow \text{Vec } A \ n \rightarrow \text{Vec } A \ n$.
- ▶ Types with modalities: $A \rightarrow_{\alpha} B$.

Types for Quantum Computing

The basic types in Quantum Computing.

- ▶ **Bit:** $|0\rangle, |1\rangle$.
- ▶ **Qubit:** $|\phi\rangle = \alpha|0\rangle + \beta|1\rangle$, where $\alpha, \beta \in \mathbb{C}, |\alpha|^2 + |\beta|^2 = 1$.
- ▶ Multi-qubits are represented by a tensor product.
Qubit \otimes Qubit, Qubit \otimes Qubit \otimes Qubit, Qubit \otimes Bit, etc.

Qubits are resource

- ▶ No cloning: *one can not duplicate a qubit.*

~~dup x = (x, x)~~

- ▶ Qubit does not exist in a vacuum.

`Init0` : Unit \multimap Qubit

let x = `Init0` () in ...

- ▶ Qubit does not disappear into the ether.

`Discard` : Qubit \multimap Unit

let x = `Init0` () in ...

let _ = `Discard` x in ...

Updating Qubits: unitary operations

One way to update qubits is via *unitary operations*.

- ▶ Reversibility: $UU^\dagger = U^\dagger U = I$.
- ▶ Linearity: $U(\alpha|0\rangle + \beta|1\rangle) = \alpha U|0\rangle + \beta U|1\rangle$.

Common quantum gates

- ▶ Hadamard gate.

$$H|0\rangle = 1/\sqrt{2}(|0\rangle + |1\rangle)$$

$$H|1\rangle = 1/\sqrt{2}(|0\rangle - |1\rangle)$$

- ▶ Phase gate.

$$S|0\rangle = |0\rangle$$

$$S|1\rangle = i|1\rangle$$

- ▶ T gate.

$$T|0\rangle = |0\rangle$$

$$T|1\rangle = \omega|1\rangle, \text{ where } \omega^2 = i$$

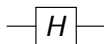
- ▶ CNOT gate.

$$\text{CNOT}|00\rangle = |00\rangle \quad \text{CNOT}|01\rangle = |01\rangle$$

$$\text{CNOT}|10\rangle = |11\rangle \quad \text{CNOT}|11\rangle = |10\rangle$$

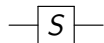
Types for quantum gates

- ▶ Hadamard gate.



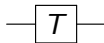
$$H : \text{Qubit} \rightarrow \text{Qubit}$$

- ▶ Phase gate.



$$S : \text{Qubit} \rightarrow \text{Qubit}$$

- ▶ T gate.



$$T : \text{Qubit} \rightarrow \text{Qubit}$$

- ▶ CNOT gate.



$$\text{CNOT} : \text{Qubit} \otimes \text{Qubit} \rightarrow \text{Qubit} \otimes \text{Qubit}$$

Measurement

Measurement is needed to readout the bit information from qubit.

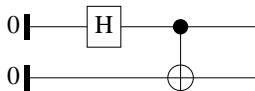


Meas : Qubit \rightarrow Bit

- ▶ $M(\alpha|0\rangle + \beta|1\rangle) = |0\rangle$ with probability $|\alpha|^2$.
- ▶ $M(\alpha|0\rangle + \beta|1\rangle) = |1\rangle$ with probability $|\beta|^2$.

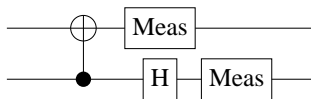
Programming quantum circuits in Proto-Quipper

```
bell00 : !(Unit -> Qubit * Qubit)
bell00 u =
  let a = Init0 ()
      b = Init0 ()
  in CNot b (H a)
```



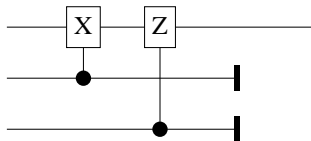
Programming quantum circuits in Proto-Quipper

```
alice : !(Qubit -> Qubit -> Bit * Bit)
alice a q =
  let (a, q) = CNot a q
      q = H q
  in (Meas a, Meas q)
```



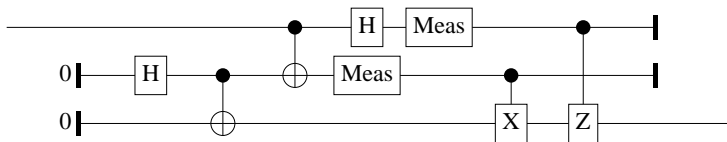
Programming quantum circuits in Proto-Quipper

```
bob : !(Qubit -> Bit -> Bit -> Qubit)
bob q x y =
  let (q, x) = C_X q x
      (q, y) = C_Z q y
      _ = Discard x
      _ = Discard y
  in q
```

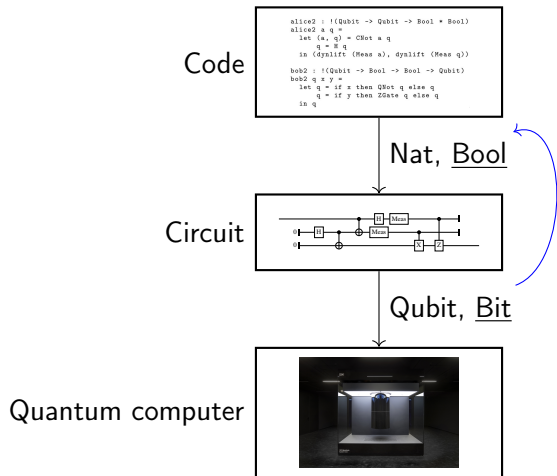


Programming quantum circuits in Proto-Quipper

```
tele : !(Qubit -> Qubit)
tele q =
  let (b, a) = bell100 ()
      (x, y) = alice a q
      z = bob b x y
  in z
```



Interleaving circuit generation time and circuit execution via dynamic lifting



Types for dynamic lifting

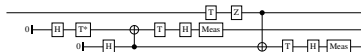
- ▶ $\Gamma \vdash_{\alpha} M : A$, where $\alpha = 0 \mid 1$.
- ▶ Dynamic lifting.

$$\frac{\Gamma \vdash_{\alpha} M : \text{Bit}}{\Gamma \vdash_0 \text{dynlift } M : \text{Bool}}$$

- ▶ Type system distinguishes computation that uses dynamic lifting vs computation that corresponds to quantum circuits.

Programming with dynamic lifting

```
v3 : !(Qubit -> Qubit)
v3 q =
  let a1 = tgate_inv (H (Init0 ()))
      a2 = H (Init0 ())
      (a1, a2) = CNot a1 a2
      a1 = H (TGate a1)
  in if dynlift (Meas a1)
     then
       let _ = Discard (Meas a2)
           in v3 q
     else let q = ZGate (TGate q)
           (a2, q) = CNot a2 q
           a2 = H (TGate a2)
           in if dynlift (Meas a2)
              then v3 (ZGate q)
              else q
```



Future research

- ▶ How do we verify the correctness of a quantum program?
 - ▶ How to prove two quantum circuits are equal?
 - ▶ How to develop tests to ensure the programs perform correctly?
- ▶ How do we compile a high-level quantum programs to lower level languages (e.g., QIR, OpenQasm)?
- ▶ Suppose we have a 127 Qubits machine, what algorithms should we run on it?

Thank you!