

781- 2013 04 02, 2013 04 04

Note Title

2013-04-02

Graded policy w/ project:

• Homework 40 → 30 → 25

Midterm 15 → 10

Final 35 → 30 → 25 → 30

Presentation 10

Project 15 → 25

Exercise 61 (Schöningh)

Convert $F = (\forall x \exists y P(x, g(y, f(x))) \vee \neg Q(z)) \vee$

$\forall x \exists u P(x, g(u, f(x))) \vee \neg Q(z) \vee \neg \forall w R(w, y)$ (rename)

$\forall x \exists u P(x, g(u, f(x))) \vee \neg Q(z) \vee \exists w \neg R(w, y)$ (push negation in)

$\forall x \exists u \exists w [P(x, g(u, f(x))) \vee \neg Q(z) \vee \neg R(x, y)]$ (prenex)

Exercise 62 (Schöningh)

Find the Skolem form of the formula

$$\forall x \exists y \forall z \exists w (\neg P(a, w) \vee Q(f(x), y))$$

$$\forall x \forall z \exists w (\neg P(a, w) \vee Q(f(x), g(x))) \quad (\text{replace } y \text{ with } g(x))$$

$$\forall x \forall z (\neg P(a, h(x, z)) \vee Q(f(x), g(x))) \quad (\text{replace } w \text{ with } h(x, z))$$

Exercise 63. Transform to rectified prenex Skolem form.

$$\forall z \exists y (P(x, g(y), z) \vee \neg \forall x Q(x)) \wedge \neg \forall z \exists x \neg R(f(x, z), z)$$

$$\forall z \exists y (P(x, g(y), z) \vee \exists x \neg Q(x)) \wedge \exists z \forall x R(f(x, z), z)$$

$$\forall z \exists y (P(x, g(y), z) \vee \exists w \neg Q(w)) \wedge \exists t \forall s R(f(s, t), t)$$

$$\forall z \exists y \exists w \exists t \forall s [P(x, g(y), z) \vee \neg Q(w) \wedge R(f(s, t), t)]$$

$$\forall z \exists w \exists t \forall s [P(x, g(\underline{h_1(z)}), z) \vee \neg Q(w) \wedge R(f(s, t), t)]$$

$$\forall z \exists t \forall s [P(x, g(\underline{h_1(z)}), z) \vee \neg Q(\underline{h_1(z)}) \wedge R(f(s, t), t)]$$

$$\forall z \forall s [P(x, g(h(z)), z) \vee \neg Q(h_1(z)) \wedge R(f(s, h_2(z)), h_2(z))]$$

Presentations

Tuesday, 4/16; Sakhib & Aikjo

Thursday, 4/18 ; Walker + Omar

Tuesday, 4/23 ; Selvi & McCosh

Example of transformation of a formula to an S -equivalent formula in rectified prenex Skolem form.

Recall,

A formula F is S -equivalent to a formula G iff F is satisfiable iff G is satisfiable.

(Motivation: equivalence requires equality of models, models are a special kind of satisfiable structures, after Skolemizing new functions

are introduced, so structures are no longer suitable. $F \Rightarrow F_{sr}$. The structures that are suitable for F are not suitable for F_{sr} , b/c F_{sr} has additional functions.)

(Example on pp. 60-61 of Schöning.)

$$F = (\neg \exists x P(x, z) \vee \forall y Q(x, f(y))) \vee \forall y P(g(x, y), z)$$

1. Rename bound variables

$$(\neg \exists x (P(x, z) \vee \forall y Q(x, f(y))) \vee \forall w P(g(x, w), z))$$

2. (Justified by exercise 48:

$F(\dots x \dots)$ is satisfiable iff $\exists x F(\dots x \dots)$ is satisfiable, where x is free in F .)

$$\exists z ((\neg \exists x (P(x, z) \vee \forall y Q(x, f(y))) \vee \forall w (P(g(x, w), z)))$$

3. Convert to prenex form

De Morgan's law as the negation is pushed in
↓

$$\exists z (\forall x (\neg P(x, z) \wedge \exists y \neg Q(x, f(y))) \vee \forall w P(g(x, w), z))$$

$$\exists z \forall x \exists y \forall w ((\neg P(x, z) \wedge \neg Q(x, f(y))) \vee P(g(x, w), z))$$

4. Skolemize

$$\exists z a \quad \forall x \exists y \forall w ((\neg P(x, a) \wedge \neg Q(x, f(y))) \vee P(g(x, w), a))$$

$$\exists y h(x) \quad \forall x \forall w ((\neg P(x, a) \wedge \neg Q(x, f(h(x)))) \vee P(g(x, w), a)) = F_y$$

5. Convert the matrix of the formula F_n into CNF.

$$F_5 = \forall x \forall w \{ (\neg P(x, a) \vee P(g(x, w), a)) \wedge (\neg Q(x, f(h(x))) \vee P(g(x, w), a)) \}$$

We can write F_5 as a clause set

$$\{ \neg P(x, a), P(g(x, w), a) \} \quad \{ \neg Q(x, f(h(x))), P(g(x, w), a) \}$$

This is the form needed for resolution refutation proofs in predicate calculus.