

We start with the (by now familiar) grammar 2.9, which is the unambiguous version of grammar 2.4

$$T \rightarrow R$$

$$T \rightarrow aTc$$

$$R \rightarrow$$

$$R \rightarrow bR$$

2.9

$$T \rightarrow R$$

$$T \rightarrow aTc$$

$$R \rightarrow$$

$$R \rightarrow RbR$$

2.4

We add a production $S' \rightarrow S$, where S is the (original) start symbol

$$0: T' \rightarrow T$$

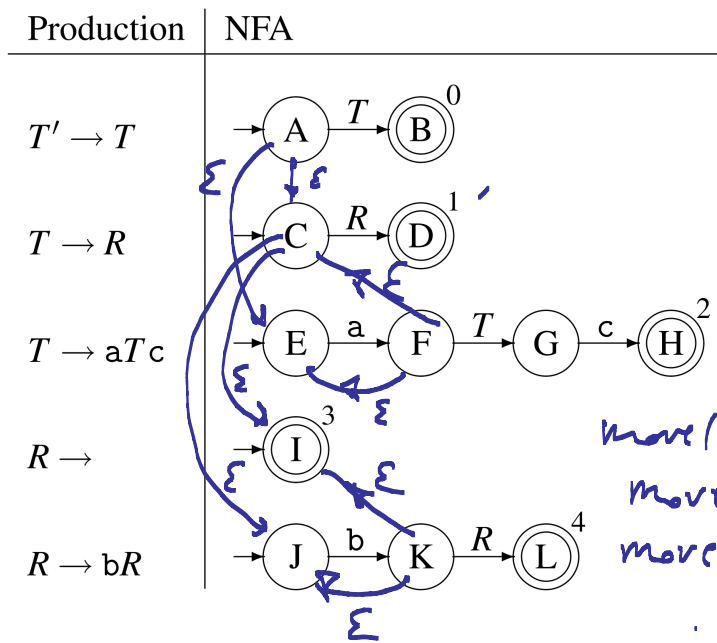
$$1: T \rightarrow R$$

$$2: T \rightarrow aTc$$

$$3: R \rightarrow$$

$$4: R \rightarrow bR$$

We convert each production to an NFA



We combine the individual NFAs for each production with ϵ -transitions linking states correspondingly to the same nonterminals:

$move(0, a) = \epsilon\text{-closure}(\{F\}) = \{C, E, F, I, J\} = 3$
 $move(0, b) = \{K, I, J\} = 4$
 $state$ | ϵ -transitions

$move(3, R) = \{D\} = 2$
 $move(3, a) = 3$ $move(3, T) = \{A\} = 5$
 $move(3, b) = 4$
 $move(4, R) = \{L\} = 6$
 $move(5, c) = \{H\} = 7$

$\epsilon\text{-closure}(\{A\}) = \{A, C, E, I, J\} = 0$
 $move(0, T) = \epsilon\text{-closure}(\{B\}) = \{B\} = 1$
 $move(0, R) = \epsilon\text{-closure}(\{D\}) = \{D\} = 2$

| state | ϵ -transitions |
|-------|-------------------------|
| A | C, E |
| C | I, J |
| F | C, E |
| K | I, J |

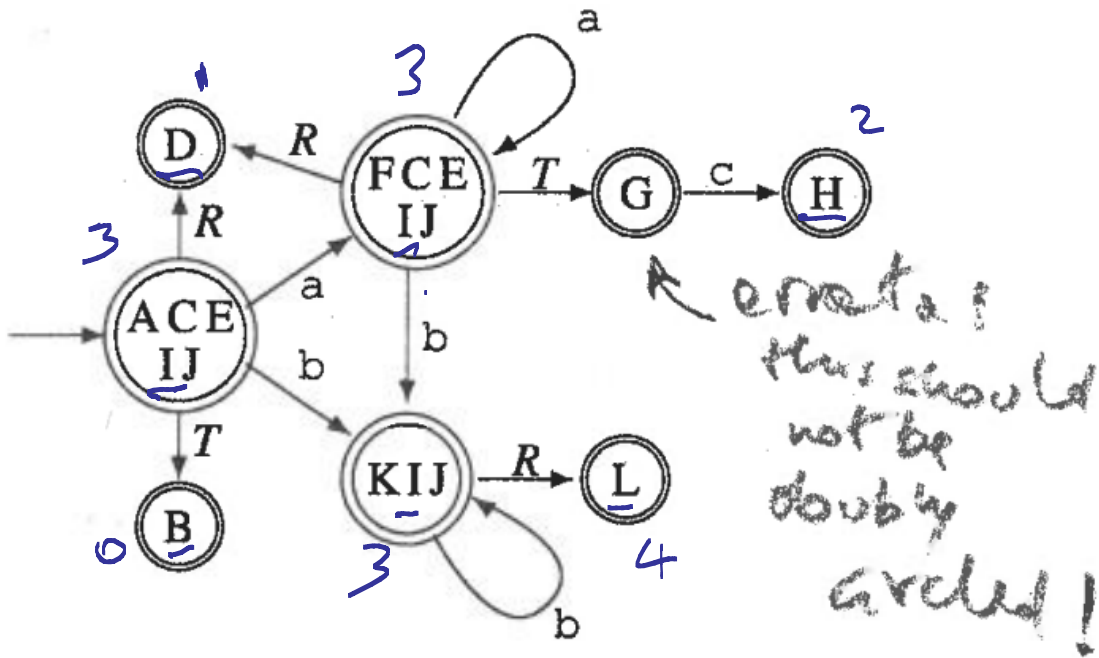


Fig. 2.32 [M]

