

CS/E 531 2020-01-16

Example of loop invariant for a program fragment to sum the integers from 1 to n

```

      i := 1;
P  -  sum := 0;
      while i ≤ n do
          sum := sum + i;
          i := i + 1;
      od;

```

while B do S od

$Q =$

$P: i=1 \wedge \text{sum}=0$ $Q: \text{sum} = \sum_{j=1}^n j$

$I: \text{sum} = \sum_{j=1}^{i-1} j \wedge i \leq n+1$

$$I \wedge \sim B \Rightarrow Q \quad \checkmark$$

$$\text{sum} = \sum_{j=1}^{i-1} j \wedge i \leq n+1 \wedge i \geq n+1 \Rightarrow \text{sum} = \sum_{j=1}^n j$$

$$\text{sum} = \sum_{j=1}^{i-1} j \wedge i = n+1$$

$$\text{sum} = \sum_{j=1}^{n+1-1} j = \sum_{j=1}^n j$$

$$P \Rightarrow I \quad \checkmark$$

$$i=1 \wedge \text{sum}=0 \Rightarrow \text{sum} = \sum_{j=1}^{i-1} j \quad \left(\begin{array}{l} \text{the upper index} \\ \text{is greater} \\ \text{than the lower one} \\ \text{when } P \text{ holds} \end{array} \right)$$

$\text{sum} = 0$

$$I \wedge B \{ SL \} I$$

If $I \wedge B$ hold, then I holds after executing SL .

I use primes ($'$) to indicate the value of a variable after executing the body of the loop.

$$\underbrace{\sum_{j=0}^{i-1} 1 \wedge i \leq n+1 \wedge i \leq n}_{\substack{I \\ B}} \Rightarrow \underbrace{\sum_{j=1}^{i'} 1 \wedge i' \leq n+1}_{I \text{ after executing } SL}$$

$$\sum_{j=0}^{i-1} 1 \wedge i \leq n$$

I

$$\text{sum}' = \text{sum} + i'$$

$$i' = i + 1 \quad i' + 1 =$$

$$\text{sum} + i' = \sum_{j=1}^{i'+1} 1 \wedge i + 1 \leq n + 1 \quad \checkmark$$

(rewrite)

$\langle \text{Start} \rangle \Rightarrow \langle \text{Letter} \rangle \Rightarrow b$

$\langle \text{Start} \rangle \Rightarrow \langle \text{start} \rangle \langle \text{Letter} \rangle \Rightarrow \langle \text{Letter} \rangle \langle \text{Letter} \rangle \Rightarrow$

$\Rightarrow b \langle \text{Letter} \rangle \Rightarrow bc$

Two derivations. The ^{string at the} end of a derivation is called a sentence. An intermediate string is called a sentential form.

The second derivation is a left-most derivation.

$\langle \text{Start} \rangle \Rightarrow \langle \text{Start} \rangle \langle \text{Letter} \rangle \Rightarrow \langle \text{Start} \rangle c \Rightarrow \langle \text{Letter} \rangle c \Rightarrow$

$\Rightarrow bc$, (A rightmost derivation)

$\langle \text{Start} \rangle \Rightarrow \langle \text{start} \rangle \langle \text{Letter} \rangle \Rightarrow \langle \text{Letter} \rangle \langle \text{Letter} \rangle \Rightarrow \langle \text{Letter} \rangle c \Rightarrow$

$\Rightarrow b c$

(neither leftmost nor rightmost)