In other words: When Dijkstra’s algorithm
from a to head (e) is a shortest path from a to head (e), is a shortest path
the path consisting of a directed to the end of a
all edges with one vertex in \( V \) and one in \( V' \), then
if \( e \) is chosen to minimize \( d(u, v) + w(e) \) over
\( V \cap V' \), \( e \in \mathcal{E} \)

\[ e = (u, v, w) \], \( (V, E, W) \). The

Thm. (C. J. Borse, 1988).

I. Correctness of Dijkstra’s algorithm.
\[ g(z^{-1}) + c(y, z) = \rho(z^{-1}, z^2) = g(z^{-1}) + c(y, z) \]

\[ g(y) + c(y, z) = (y, z) \]

\[ w(p) = \rho(y, y) + w(e) \]

From \( y \) to \( z \).

Be any other path.

Let \( p = r_1, z^e, r_2, \ldots \), \( p = r_1, y, r_2, \ldots \).

Choose \( e \) and

\( y \rightarrow z \)

Proof:
Proof. By an inductive argument, we prove that the shortest path from each node to the source has a length of $f(p) = \text{length of shortest path from } p$ to the source. Then solve for the single-source, single-directed shortest paths using any of the Dijkstra's shortest path algorithms. The least number of edges among all paths that reach the source is the desired shortest path.
problem solved by Bigtooth's algo n' em.

something can be drawn from the leading to the

ordered by non-decreasing desceent from 5

due to C and a sequence of all variables in G

@. last non-adjacent produces a shortest-path spanning

and consistent with &apos;soooooo Happy. &apos; (n log n + m)

strong aptitude for O.K. they're a whole bunch when