

CSCE 747 Software Testing and Quality Assurance

Lecture 07 – Dataflow Testing

Last Time

- Lecture 06 Slides 1-19 covered last time
- Case Study Question after class
- Path Testing continued
- Ch 9 pp 131-149

Today

- Dataflow Testing
- Ch 10, pp 151-167

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Dataflow Testing

- Dataflow testing refers to forms of structural testing that focus on:
 - the points at which variables ^{Def} receive values and
 - the points at which these values are used
- "Dataflow testing serves as a reality check on path testing;"

Forms of dataflow testing

- **Two main forms of dataflow testing:**
 - One provides a set of basic definitions and a unifying structure of test coverage metrics
 - The other based on a concept called a program slice.
- **Start with program graph but move back towards functional testing**

Define/Use of Variables

- Define/Use information
 - $x = y + 3 * z$
 - Define a new x; use variables y and z
- Concordances that list statement numbers in which variable names occur *x*
- Define/reference anomalies:
 - A variable that is defined but never used
 - A variable that is used before it is defined
 - A variable that is defined twice before it is used

Static Analysis

- Static analysis: finding faults in source code without executing it.

Define/Use Testing

- **define/use testing was done by**
 - Rapps and Weyuker, IEEE Transactions on Software Engineering, Vol. SE-11, 1985
- **Definitions: Given a program P**
 - $G(P)$ – the program graph; single entry; single exit
 - $PATHS(P)$ – the set of all paths in P

“Definition” and “Usage” nodes for variable

- Definition: Node $n \in G(P)$ is a defining node of the variable $v \in V$, written as DEF(v, n)
- Definition: Node $n \in G(P)$ is a usage node of the variable $v \in V$, written as USE(v, n)
 - ✗ Definitions(n) – variables v that are defined in statement n
 - Usage(n) – variables that are used in statement n
 - ■ Definitions(v) – statements that define v
 - ■ Usage(v) – statements that use v
 - Next-Use(n, v) – list of statements following n that use v
- Node n : statement fragment $x = y + z$

Def(x, n)

Example (Commission)

10. totalLocks = 0 Defs (t, L) = {10, 16}

11. totalStocks = 0

12. totalBarrels = 0

13. Input(locks)

14. While NOT(locks = -1) 'loop condition uses -1

15. Input(stocks, barrels)

16. totalLocks = totalLocks + locks

17. totalStocks = totalStocks + stocks

18. totalBarrels = totalBarrels + barrels

19. Input(locks)

20. EndWhile

21. Output("Locks sold: ", totalLocks)

Predicate/Computation Use

- $USE(v, n)$ can be classified as
 - Predicate use (P-use)
 - Computation use (C-use)
- $USE(a, 7)$?
- $USE(a, 9)$?

'Step 2: Is A Triangle? Modified

6. $t1 = b + c$

7. $t2 = \underline{a} + c$

8. $t3 = a + b$

9. If ($\underline{a} < t1$) AND ($b < t2$) AND ($c < t3$)

10. Then IsATriangle = True

11. Else IsATriangle = False

12. EndIf

...

definition-use path

- Definition: A definition-use path with respect to a variable v (denoted $\text{du-path}(v)$) is a path in $\text{PATHS}(P)$ such that
- ~~for some~~ $v \in V$, there are define and usage nodes $\text{DEF}(v, m)$ and $\text{USE}(v, n)$ such that
- m and n are the initial and final nodes of the path.

Definition-Clear Path

- Definition: A definition-clear path with respect to a variable v (denoted dc-path) is a definition-use path in $\text{PATHS}(P)$
- with initial and final nodes $\text{DEF}(v, m)$ and $\text{USE}(v, n)$ such that no other node in the path is a defining node of v


The diagram shows a horizontal line representing a path. On the left, the word 'Def' is written in red. In the middle, the word 'DEF' is written in red and crossed out with a red 'X'. On the right, the word 'USE' is written in red. A red line connects 'Def' to 'USE'.
- Du-paths and dc-paths describe the flow of data across source statements from points at which the values are defined to points at which the values are used.
- Du-paths that are not definition-clear are potential trouble spots.

Compilers Again: Register Allocation

1. Program Commission (INPUT,OUTPUT)

2. Dim ...

7. lockPrice = 45.0

8. stockPrice = 30.0

9. barrelPrice = 25.0

10. totalLocks = 0

11. totalStocks = 0

12. totalBarrels = 0

13. Input(locks)

14. While NOT(locks = -1)

15. Input(stocks, brrls)

16. totalLocks = totalLocks + locks

17. totalStocks = totalStocks + stocks

18. totalBarrels = totalBarrels + brrls

19. Input(locks)

20. EndWhile

21. Output("Locks sold: ", totalLocks)

22. Output("Stocks sold: ", totalStocks)

23. Output("Barrels sold: ", totalBarrels)

24. lockSales = lockPrice * totalLocks — USE

25. stockSales = stockPrice * totalStocks

26. barrelSales = barrelPrice * totalBarrels

27. sales = lockSales + stockSales + barrelSales

28. Output("Total sales: ", sales)

29. If (sales > 1800.0)

30. Then

31. commission = 0.10 * 1000.0

32. commission = comm. + 0.15 * 800.0

33. comm. = comm. + .20 * (sales - 1800.0)

34. Else If (sales > 1000.0)

35. Then

36. commission = 0.10 * 1000.0

37. comm = comm + .15 * (sales - 1000)

38. Else

39. commission = 0.10 * sales

40. EndIf

41. EndIf

42. Output("Commission is \$" commission)

43. End Commission

*S (locks 16)
= 13, 19, 14, 20.*

DD-Paths in Figure 10.1 (previous slide)

- Table 10.1 DD-Paths in Figure 10.1

DD-Path

Nodes

A

7, 8, 9, 10, 11, 12, 13

~~X~~ Not du-path

B

14

C

15, 16, 17, 18, 19, 20 — loop

D

21, 22, 23, 24, 25, 26, 27, 28

E

29

F

30, 31, 32, 33

G

34

H

35, 36, 37

I

38, 39

J

40

K

41, 42, 43

Fig 10.1
program graph
for commission
program

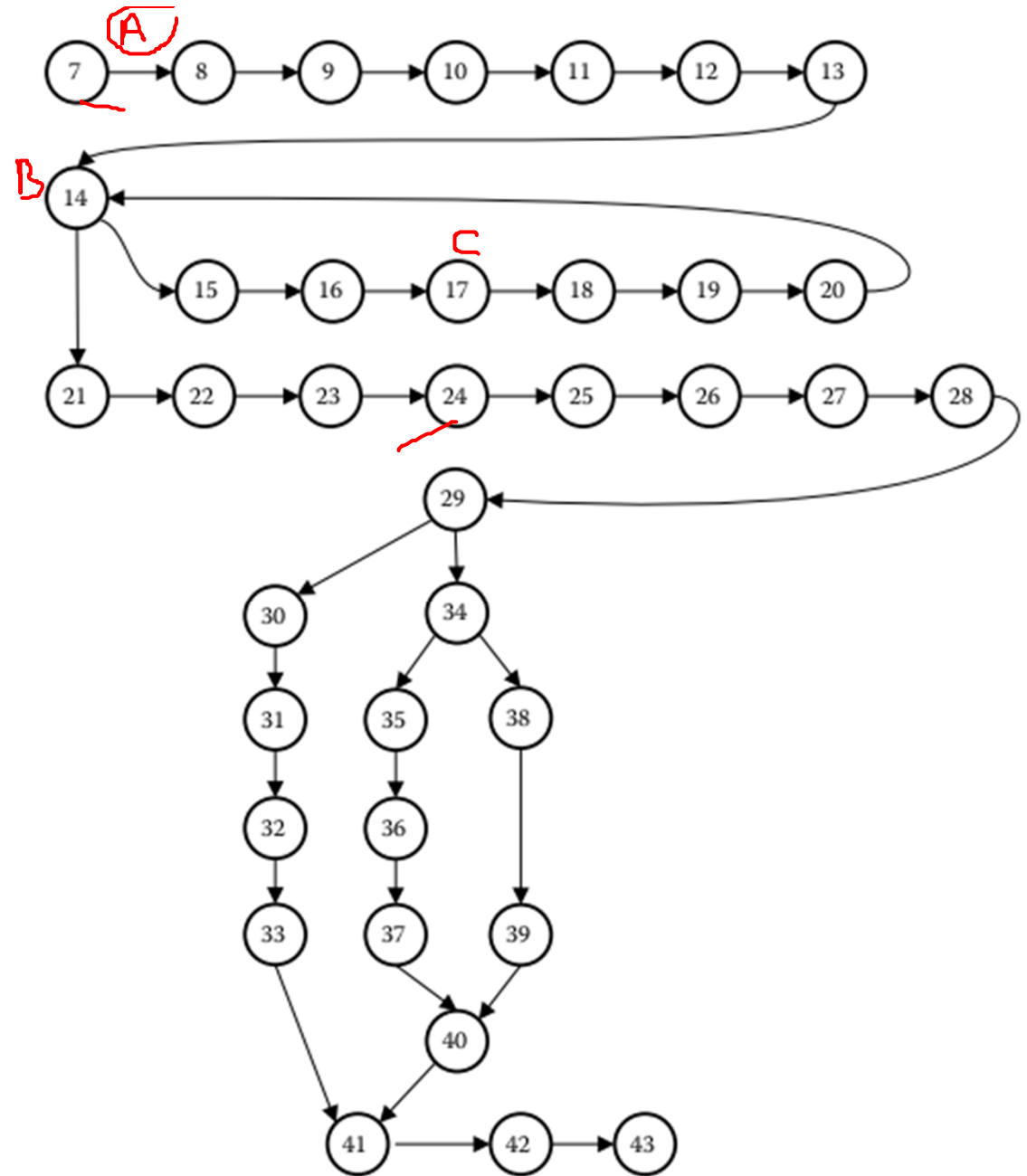
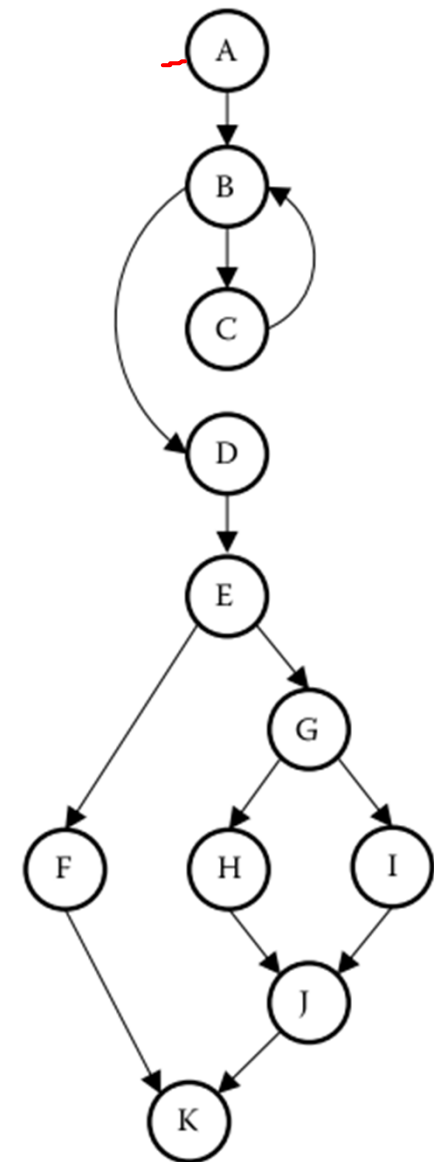


Figure 10.2 DD- Path graph of the commission program.



10.1.2 du-Paths for Stocks

- Def(stocks, 15)
- Use(stocks, 17)

10.1.3 du-Paths for Locks

- p1 = <13, 14> exit
 - p2 = <13, 14, 15, 16>
 - p3 = <19, 20, 14>
 - p4 = <19, 20, 14, 15, 16>
- 

10.1.4 du-Paths for Total-Locks

- $p6 = \langle 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 14, 21 \rangle$
- $p7 = \langle 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 14, 21, 22, 23, 24 \rangle$
- $p7 = \langle p6, 22, 23, 24 \rangle$
- $p8 = \langle 16, 17, 18, 19, 20, 14, 21 \rangle$
- $p9 = \langle 16, 17, 18, 19, 20, 14, 21, 22, 23, 24 \rangle$

$p_x = 16 \text{ ————— } 14 \text{ ————— } 16$

Table 10.2 Define/Use Nodes for Variables in the Commission Problem

12efslv

<i>Variable</i>	<i>Defined at Node</i>	<i>Used at Node</i>
lockPrice	7	24
stockPrice	8	25
barrelPrice	9	26
totalLocks	10, 16	16, 21, 24
totalStocks	11, 17	17, 22, 25
totalBarrels	12, 18	18, 23, 26
locks	13, 19	14, 16
stocks	15	17
barrels	15	18
lockSales	24	27
stockSales	25	27
barrelSales	26	27
sales	27	28, 29, 33, 34, 37, 39
commission	31, 32, 33, 36, 37, 39	32, 33, 37, 42

10.1.5 du-Paths for Sales

- $p_{10} = \langle 27, 28 \rangle$
- $p_{11} = \langle 27, 28, 29 \rangle$
- $p_{12} = \langle 27, 28, 29, 30, 31, 32, 33 \rangle$
- $p_{13} = \langle 27, 28, 29, 34 \rangle$
- $p_{14} = \langle 27, 28, 29, 34, 35, 36, 37 \rangle$
- $p_{15} = \langle 27, 28, 29, 34, 38, 39 \rangle$

Du-Paths for commission

Table 10.3 Selected Define/Use Paths

<i>Variable</i>	<i>Path (Beginning, End)</i>	
	<i>Nodes</i>	<i>Definition-Clear?</i>
lockPrice	7, 24	Yes
stockPrice	8, 25	Yes
barrelPrice	9, 26	Yes
totalStocks	11, 17	Yes
totalStocks	11, 22	No
totalStocks	11, 25	No
totalStocks	17, 17	Yes
totalStocks	17, 22	No
totalStocks	17, 25	No
locks	13, 14	Yes
locks	13, 16	Yes
locks	19, 14	Yes
locks	19, 16	Yes
sales	27, 28	Yes
sales	27, 29	Yes
sales	27, 33	Yes
sales	27, 34	Yes
sales	27, 37	Yes
sales	27, 39	Yes

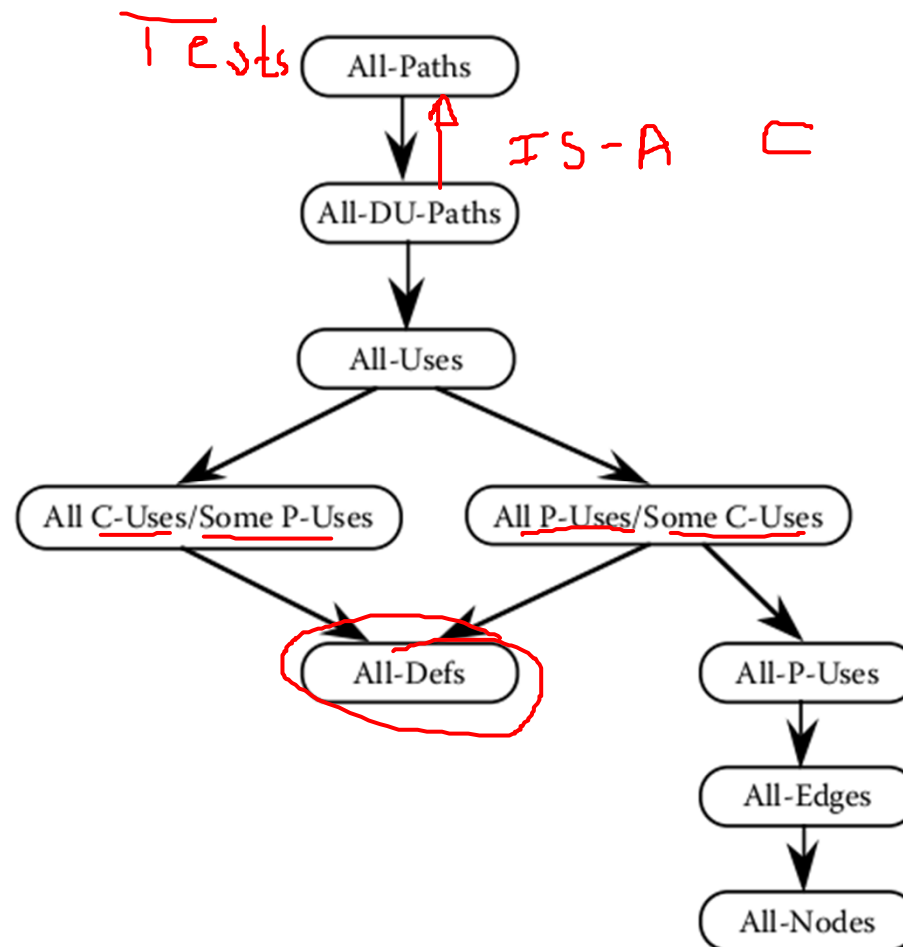
Table 10.4 Define/Use Paths for Commission

<i>Variable</i>	<i>Path (Beginning, End) Nodes</i>	<i>Feasible?</i>	<i>Definition-Clear?</i>
commission	31, 32	Yes	Yes
commission	31, 33	Yes	No
commission	31, 37	No	n/a
commission	31, 42	Yes	No
commission	32, 32	Yes	Yes
commission	32, 33	Yes	Yes
commission	32, 37	No	n/a
commission	32, 42	Yes	No
commission	33, 32	No	n/a
commission	33, 33	Yes	Yes
commission	33, 37	No	n/a
commission	33, 42	Yes	Yes
commission	36, 32	No	n/a
commission	36, 33	No	n/a
commission	36, 37	Yes	Yes
commission	36, 42	Yes	No
commission	37, 32	No	n/a
commission	37, 33	No	n/a
commission	37, 37	Yes	Yes
commission	37, 42	Yes	Yes
commission	38, 32	No	n/a
commission	38, 33	No	n/a
commission	38, 37	No	n/a
commission	38, 42	Yes	Yes

10.1.7 du-Path Test Coverage Metrics

- Definition: The set T satisfies the **All-Defs criterion** for the program P iff for every variable $v \in V$, T contains definition-clear paths from every defining node of v to a use of v .
- Definition: The set T satisfies the **All-Uses criterion** for the program P iff for every variable $v \in V$, T contains definition-clear paths from every defining node of v to every use of v , and to the successor node of each $USE(v, n)$.
- Definition: The set T satisfies the **All-P-Uses/Some C-Uses criterion** for the program P iff for every variable $v \in V$, T contains definition-clear paths from every defining node of v to every predicate use of v ;
- if a definition of v has no P-uses, a definition-clear path leads to at least one computation use.

Rapps–Weyuker hierarchy of dataflow coverage metrics.





- **Definition:** The set T satisfies the All-C-Uses/Some P-Uses criterion for the program P iff for every variable $v \in V$, T contains definition-clear paths from every defining node of v to every computation use of v ;
- if a definition of v has no C-uses, a definition-clear path leads to at least one predicate use.
- **Definition:** The set T satisfies the All-du-paths criterion for the program P iff for every variable $v \in V$, T contains definition-clear paths from every defining node of v to every use of v and to the successor node of each $USE(v, n)$, and that these paths are either single-loop traversals or cycle-free.

Slices



- **Slices of History:**

- American History
- Western Civ.
- Chinese History
- Russian History
- ... history focusing on one group

HIST 201, 202
HIST 101, 102

- **Slices of program**

- “set of program statements that contributes to or
- affects a value for a variable at some point in the program”

10.2 Slice-Based Testing

- Program slices have surfaced and submerged in software engineering literature since the early 1980s.
- They were originally proposed in Weiser (1988), used as an approach to software maintenance in Gallagher and Lyle (1991), and more recently
- used to quantify functional cohesion in Bieman and Ott (1994)
- ^vPart of this versatility is due to the natural, intuitively clear intent of the program slice concept.^v
- Informally, a program slice is a set of program statements that contributes to or affects a value for a variable at some point in the program.

slice on the variable set V

- Definition: Given a program P and a set V of variables in P , a slice on the variable set V at statement n , written $S(V, n)$, is the set of all statements in P prior to node n that contribute to the values of variables in V at node n .

$USE(V, n)$

Use nodes

- USE relationship pertains to five forms of usage:
- P-use Used in a predicate (decision) ✓
- C-use Used in computation ✓
- O-use Used for output ✓
- L-use Used for location (pointers, subscripts) ✓
- I-use Iteration (internal counters, loop indices)

Definition Nodes

- two forms of definition nodes:
- I-def Defined by input
- A-def Defined by assignment

- Slices on the locks

- S1: $S(\text{locks}, 13) = \{13\}$

- S2: $S(\text{locks}, 14) = \{13, 14, 19, 20\}$

- S3: $S(\text{locks}, 16) = \{13, 14, 19, 20\}$

- Jorgensen, Paul C. (2011-07-16). Software Testing (Page 162). Auerbach Publications. Kindle Edition.

Slices and Loops

- S9: $S(\text{totalLocks}, 10) = \{10\}$
- S10: $S(\text{totalLocks}, 16) = \{10, 13, 14, 16, 19, 20\}$
- S11: $S(\text{totalLocks}, 21) = \{10, 13, 14, 16, 19, 20\}$

Handwritten annotations for S11:
- A red circle around the value 10, with the word "init 0" written below it.
- A red circle around the value 20, with the word "locks" written below it.
- A red line under the entire set notation $\{10, 13, 14, 16, 19, 20\}$, with the phrase "total locks for lock" written below it.
- Red arrows point from the words "locks" and "lock" to the values 13 and 14 respectively.

- S12: $S(\text{totalStocks}, 11) = \{11\}$
- S13: $S(\text{totalStocks}, 17) = \{11, 13, 14, 15, 17, 19, 20\}$
- S14: $S(\text{totalStocks}, 22) = \{11, 13, 14, 15, 17, 19, 20\}$
- S15: $S(\text{totalBarrels}, 12) = \{12\}$
- S16: $S(\text{totalBarrels}, 18) = \{12, 13, 14, 15, 18, 19, 20\}$
- S17: $S(\text{totalBarrels}, 23) = \{12, 13, 14, 15, 18, 19, 20\}$

Lattice of Slices for Commission ~~Problem~~ ^{Variant}

- S31: $S(\text{commission}, 31) = \{31\}$
- S32: $S(\text{commission}, 32) = \{31, 32\}$
- S33: $S(\text{commission}, 33) = \{7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 24, 25, 26, 27, 29, 30, 31, 32, 33\}$
- S34: $S(\text{commission}, 36) = \{36\}$
- S35: $S(\text{commission}, 37) = \{7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 24, 25, 26, 27, 36, 37\}$
- S36: $S(\text{commission}, 39) = \{7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 24, 25, 26, 27, 29, 34, 38, 39\}$
- S37: $S(\text{commission}, 41) = \{7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 24, 25, 26, 27, 29, 30, 31, 32, 33, 34, 35, 36, 37, 38, 39\}$

