

Majority-of-3 code

Alice:

$$0 \mapsto 000 = 0_L$$

$$1 \mapsto 111 = 1_L$$

Bob: majority(x,y,z)
where x,y,z are what Bob receives.

Binary symmetric error channel
with error prob $p \leq \frac{1}{2}$

physical

$$0 \mapsto 0 \text{ w/ prob } 1-p$$

$$0 \mapsto 1 \text{ " " " } p$$

$$1 \mapsto 0 \text{ " " " } p$$

$$1 \mapsto 1 \text{ " " " } 1-p$$

000 $\xrightarrow{p(1-p)^2}$ 100 or $\xrightarrow{p(1-p)^2}$ 010 or $\xrightarrow{p(1-p)^2}$ 001
or $\xrightarrow{p^2}$ 000 or $\xrightarrow{p^2}$ 111

Bob can correct

111 (similar)

\therefore Bob can correct with prob

$$(1-p)^3 + 3p(1-p)^2$$

$$= 1 - 3p + 3p^2 - p^3 + 3p - 6p^2 + 3p^3$$

$$= 1 - (3p^2 - 2p^3)$$

prob that Bob does not correct = $3p^2 - 2p^3$

$q = O(p^2)$
as $p \rightarrow 0$

Quantum bit-flip channel:
with error prob p

$$\mathcal{E}(p) = (1-p)I + pX$$

Kraus operators are

$$\{\sqrt{1-p}I, \sqrt{p}X\}$$

Alice: $|\psi\rangle := \alpha|0\rangle + \beta|1\rangle$
 $\mapsto |\psi_L\rangle := \alpha|0_L\rangle + \beta|1_L\rangle$
 $= \alpha|000\rangle + \beta|111\rangle$

sends to Bob, but Bob can't just measure the 3 qubit values because that will collapse the state, destroying the superposition.
Want Bob to recover $|\psi_L\rangle$

Idea: Bob measures the error syndrome, that is, what type of error probably occurred in the channel without getting any classical info about $|\psi_L\rangle$.

Alice: $P := |\psi\rangle\langle\psi| \in \mathcal{L}(\mathbb{C}^2)$

P travels $\xrightarrow{\mathcal{E}^{\otimes 3}}$ \mathcal{O} (recovery)

$|\psi_L\rangle\langle\psi_L|$ $\xrightarrow{\mathcal{E}^{\otimes 3}}$ \mathcal{O} (recovery)

Each qubit Bob errors on each qubit independently

Encoding (Alice)

$|\psi_L\rangle = \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$

Bob gets \mathcal{O} , some 3-qubit state.

1) measures the error syndrome by computing the parity of adjacent qubits.
Projective measurement with

$P_0 := (|000\rangle + |111\rangle)(\langle 000| + \langle 111|)$

$P_1 := (|100\rangle + |011\rangle)(\langle 100| + \langle 011|)$

$P_2 := (|101\rangle + |110\rangle)(\langle 101| + \langle 110|)$

Bob error recovery channel:

$b_1, b_2 \in \{0, 1\}$
 $b_1 = b_2 = 0 : P_0$
 $b_1 = 0, b_2 = 1 : P_1$
 $b_1 = 1, b_2 = 0 : P_2$
 $b_1 = b_2 = 1 : P_3$

Bob's error-recovery channel \mathcal{R} :

$$\mathcal{R}(\sigma) = P_0 \sigma P_0 + X_1 P_1 \sigma P_1 X_1 + X_2 P_2 \sigma P_2 X_2 + X_3 P_3 \sigma P_3 X_3$$

Kraus operators are $\{P_0, X_1 P_1, X_2 P_2, X_3 P_3\}$

Phase-flip channel (error prob p)

$$\tilde{\mathcal{E}}(p) = (1-p)P + pZPZ$$

Notice:

$$\tilde{\mathcal{E}}(H \rho H) = (1-p)H \rho H + pZ H \rho H Z$$

$$[HZ = XH, HX = ZH]$$

$$= (1-p)H \rho H + pHX \rho XH = H((1-p)\rho + pX \rho X)H = H \tilde{\mathcal{E}}(p)H$$

Alice's encoding (for phase-flip channel):

Error recovery (Bob)

Bit & phase flip:

$$\mathcal{E}(p) = (1-p)\rho + Z X \rho X Z = (1-p)\rho + Y \rho Y$$

Partially degrading channel includes all 3 possibilities (with equal prob):

$$\mathcal{D}(p) = (1-p)\rho + \frac{p}{3}(X \rho X + Y \rho Y + Z \rho Z)$$

Shor code (9 qubits)

Thm: (Discretization of errors)

If you can recover from a channel with Kraus ops K_1, \dots, K_n , then same code/recovery recovers from any channel whose Kraus ops are linear combinations of K_1, \dots, K_n

Cor: Shor code recovers from arbitrary 1-qubit errors.