

Mergeable Heaps (m-heaps)

Basic Ops

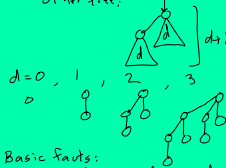
- Insert(H, x)
- FindMin(H)
- DeleteMin(H)
- DecreaseKey(H, x, k)
- Delete(H, x)
- Merge(H_1, H_2) — combine heaps H_1 & H_2 into a single heap (H_1 & H_2 are destroyed)

2 implementations:

- Binomial Heaps
- Fibonacci Heaps

A binomial tree of degree $d \geq 0$ is defined inductively as follows:

1. A binomial tree of degree 0 is a single node (the root)
2. For $d \geq 0$, a binomial tree of degree $d+1$ is obtained from 2 binomial trees of degree d by making the root of one tree the leftmost child of the root of the other tree.



Basic facts:

- A binomial tree of degree d has height d & size 2^d (by induction on d)
- Every node in a binomial tree is the root of a binomial (sub)tree
- The root of a binomial tree of degree d has exactly d many children (induction on d) (degree of a node = # of its children)
- The children of a node of degree d have degrees $d-1, d-2, \dots, 0$ from left to right

All easily provable by induction on d .

For $0 \leq i \leq d$, the i th level of a binomial tree of degree d has exactly $\binom{d}{i}$ many nodes, where $\binom{d}{i} = \frac{d!}{i!(d-i)!}$ is the binomial coefficient "d choose i"

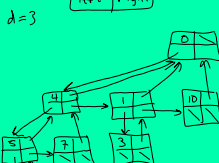
[induction on i using the "Pascal's triangle" recurrence:
 $\binom{d}{0} = \binom{d}{d} = 1$ (boundary conditions)
 for $0 < i < d$

$$\binom{d}{i} = \binom{d-1}{i-1} + \binom{d-1}{i}$$

A binomial heap is a sequence of binomial trees of strictly increasing degree.

data (key, satellite info) are kept in the nodes. Each tree in the sequence has keys in min-heap order. Nodes have the following attributes:

- data — as above
- leftmost child (left) — pointer to the leftmost child of the node
- right sibling (right) — pointer to the nodes immediate right sibling (on the same level)
- parent — points to the parent
- degree — the degree of the node



Note the shape of a binomial heap is uniquely determined by its size.