

Sums

Def: Let $a_1, a_2, \dots, a_n, \dots$ be numbers

$$\sum_{i=1}^u a_i = a_1 + a_2 + \dots + a_u$$

If $u < \infty$, then this sum is 0. (empty sum)

$$\sum_{i=1}^{\infty} a_i := \lim_{n \rightarrow \infty} \sum_{i=1}^n a_i$$

Common sums:

$$\sum_{i=1}^u 1 = u - 1 + 1 \quad (if \ u \geq 1)$$

Arithmetic sums

$$\sum_{i=1}^n i = \frac{n(n+1)}{2} = \Theta(n^2)$$

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6} = \Theta(n^3)$$

$$\sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4} = \Theta(n^4)$$

generally, (k constant), $k \geq 0$

$$\sum_{i=1}^n i^k = \Theta(n^{k+1})$$

Proof: UTS $\Theta(n^{k+1})$ and upper bound

lower bound $\Omega(n^{k+1})$

U.B: $\sum_{i=1}^n i^k \leq \sum_{i=1}^n n^k \quad (if \ i \leq n)$

$$= n^k \sum_{i=1}^n 1 = n^k \cdot n = n^{k+1}$$

\therefore U.B. (constant can be 1)

L.B:

$$\sum_{i=1}^n i^k \geq \sum_{i=\frac{n}{2}}^n i^k \quad (\text{truncation of the sum})$$

$$\geq \sum_{i=\frac{n}{2}}^n \left(\frac{n}{2}\right)^k \quad (\text{replace each term with smallest term})$$

$$\geq \left(\frac{n}{2}\right)^k \sum_{i=\frac{n}{2}}^n 1 = \frac{1}{2} n^{k+1}$$

$$= \Omega(n^{k+1}) \quad (\text{can choose } C = \frac{1}{2})$$

for $i=1$ to n sum to $\Theta(i)$

$$\sum_{i=1}^n \Theta(i) = \Theta\left(\sum_{i=1}^n i\right) = \Theta(n^2)$$

Geometric Sums

Finite geometric sum is of

$$\sum_{i=0}^{n-1} r^i$$

$r \neq 1$ & is constant "ratio"

Proof: Let $S = \sum_{i=0}^{n-1} r^i$

$$rS = r + r^2 + \dots + r^n$$

$$rS - S = r^n - 1$$

$$(r-1)S = r^n - 1 \quad \therefore S = \frac{r^n - 1}{r-1}$$

Infinite geometric sum $\left(\frac{1}{1-r}\right)$ ($|r| < 1$):

$$\sum_{i=0}^{\infty} r^i = \lim_{n \rightarrow \infty} \sum_{i=0}^n r^i = \lim_{n \rightarrow \infty} \frac{1-r^{n+1}}{1-r}$$

$$= \frac{1}{1-r} - \frac{r^{n+1}}{1-r} \xrightarrow[n \rightarrow \infty]{=0} \frac{1}{1-r}$$

for $(i=1, i < n, i \times 2)$

$\Theta(i) \left[\begin{array}{l} \text{for } (j=1, j \leq i; j++) \\ \text{sum}++; \end{array} \right.$

bth iteration of the outer loop, $i=2^{b-1}$

iterations $\approx k$ such that $2^{k-1} = n$

solve for k:

$$(2^{k-1} = n) \text{ log of both sides}$$

$$\frac{(k-1) \log 2}{\text{power rule for logs}} = \log n$$

$$k-1 = \log n$$

$$k = \log n + 1$$

Total time is $\sum_{b=1}^{\log n + 1} \Theta(2^{b-1})$

$= \Theta\left(\sum_{b=1}^{\log n + 1} 2^{b-1}\right)$ an $\log n + 1$ iterations

$= \Theta\left(\sum_{m=0}^{\log n} 2^m\right)$ $m=b-1$

$= \Theta\left(\frac{2^{\log n + 1} - 1}{2 - 1}\right)$

$= \Theta(2^{\log n + 1} - 1)$

$= \Theta(2^{\log n + 1}) = \Theta(2 \cdot 2^{\log n})$

$= \Theta(2n) = \Theta(n)$

Harmonic sum

$$H(n) = \sum_{i=1}^n \frac{1}{i} = \Theta(\log n)$$

Proof: UB:

$$H(n) = \frac{1}{B_0} + \frac{1}{B_1} + \frac{1}{B_2} + \dots + \frac{1}{B_{\log n}}$$

$$B_i = \frac{1}{2^i} + \dots + \frac{1}{2^{i-1}} \leq 2^i \cdot \frac{1}{2^i}$$

terms in B_i is i largest term in B_i

$= 1$

$B_{\lceil \log n \rceil}$ has $2^{\lceil \log n \rceil} > 2^{\log n} = n$ terms

so blocks don't go beyond $B_{\lceil \log n \rceil}$

\therefore There are $\leq 1 + \lceil \log n \rceil \leq \log n$ many blocks

So $H(n) \leq \sum_{i=0}^{\lceil \log n \rceil} B_i \leq \sum_{i=0}^{\lceil \log n \rceil} 1$

$= 1 + \lceil \log n \rceil \leq 2 + \log n$

$= O(\log n)$

L.B. similar block decomp:

$$B_i = \frac{1}{2^i} + \dots + \frac{1}{2^{i-1}}$$

$$> 2^i \cdot \left(\frac{1}{2^{2i}}\right) = \frac{1}{2}$$

blocks: last full block starts with $\frac{1}{2^k}$ for largest

k such that $\frac{1}{2^{k+1}} \geq \frac{1}{n}$

Solve for k:

$$2^{k+1} - 1 \leq n$$

$$2^{k+1} \leq n+1$$

$$k+1 \leq \log(n+1)$$

$$k \leq \log(n+1) - 1$$

k is largest such, so

$$k = \lfloor \log(n+1) - 1 \rfloor$$

$$\geq \log n - 2$$

So

Then

$$\underbrace{\int_{l-1}^u f(x) dx}_{L.B} \leq \sum_{i=l}^u f(i) \leq \underbrace{\int_l^{u+1} f(x) dx}_{U.B}$$

$$U.B - L.B = \int_u^{u+1} f(x) dx - \int_{l-1}^l f(x) dx$$

$$\leq f(u+1) - f(l-1)$$

If f is monotone descending then

$$\int_l^{u+1} f(x) dx \leq \sum_{i=l}^u f(i) \leq \int_{l-1}^u f(x) dx$$

$$\int_1^{n+1} \frac{dx}{x} \leq \sum_{i=1}^n \frac{1}{i} = 1 + \sum_{i=2}^n \frac{1}{i}$$

$$\ln(n+1) = \Theta(\lg(n+1))$$

$$= \Theta(\lg n)$$

$$\leq 1 + \int_1^n \frac{dx}{x} = 1 + \ln n = \Theta(\lg n)$$

Mixed arithmetic/geometric series

$$\sum_{i=1}^n i \cdot r^i$$