

CSCE 355
1/29/2024

DFA minimization

①

Given a DFA $A = \langle Q, \Sigma, \delta, q_0, F \rangle$,
goal is to find a DFA A_{\min} such that
 $L(A_{\min}) = L(A)$ and A_{\min} has the fewest
possible states. (A_{\min} is unique given A).

2 steps:

1) Remove from A all states not reachable
from q_0 . (A DFA where all states are
reachable we'll call sane.)
Such a DFA is equivalent to A .

2) Merge clusters of indistinguishable
states into single state.

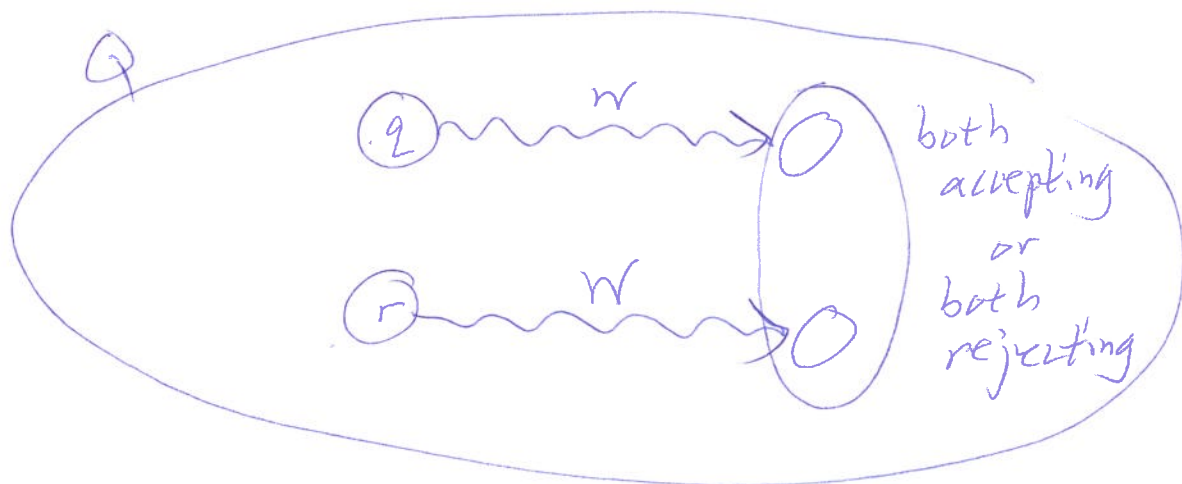
Definition: Let A_{sane} be as above and let $q \in Q$
be a state of A . Define the DFA

$$A_q := \langle Q, \Sigma, \delta, \underset{\substack{\text{start} \\ \text{state}}}{q}, F \rangle$$

[Note. $A_{q_0} = A$.]

Def: $q, r \in Q$ are indistinguishable
if $L(A_q) = L(A_r)$.

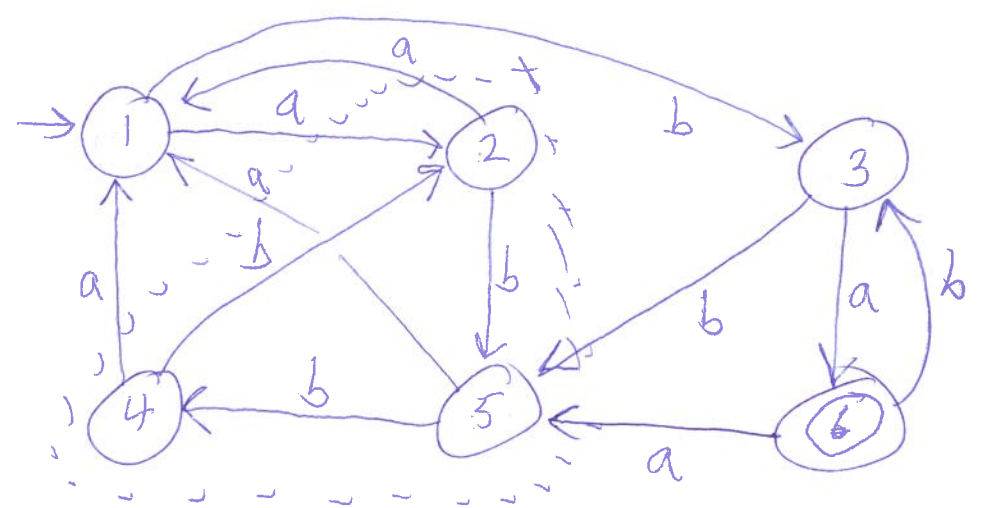
(2)



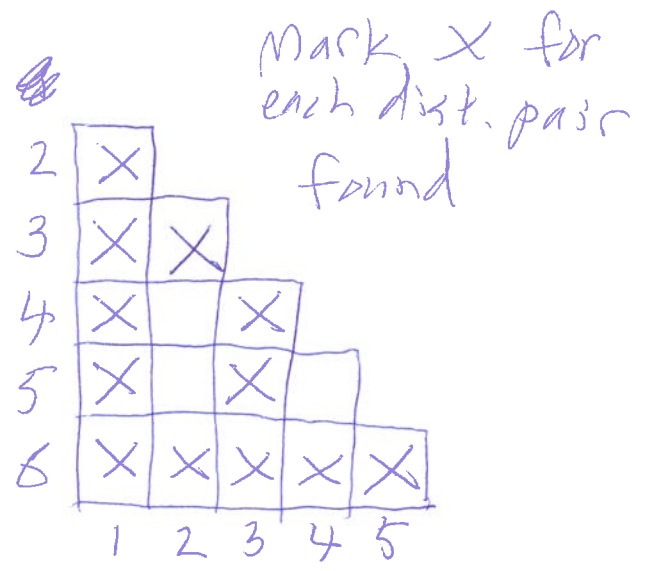
So q, r are distinguishable if $L(A_q) \neq L(A_r)$,
that is, there exists a string $w \in \Sigma^*$
such that, reading w from q gives a different
result from reading w from r . (One accepting,
the other rejecting.) w distinguishes q from r

To perform Step 2, find pairs of distinguishable
states, until we can't find any more. Pairs left over
are guaranteed indistinguishable.

Ex:
 $\Sigma = \{a, b\}$



Rule 1: Every accepting state is distinguishable from every rejecting state, (distinguished by $w = \epsilon$)



Rule 2: q & r are distinguishable if there exists $s \in \Sigma^*$ such that $\delta(q, s)$ & $\delta(r, s)$ are distinguishable.



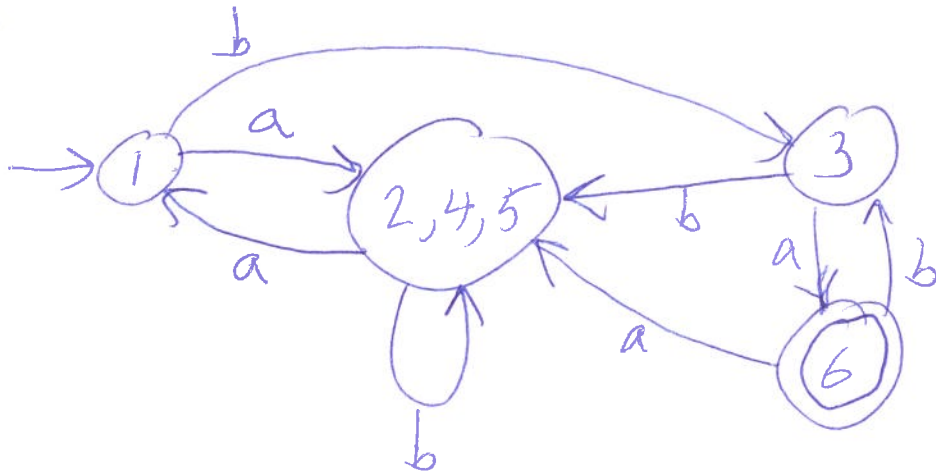
if (t, u) are dist.
 then (q, r) are dist.

Fact: Rules 1 & 2 suffice for finding all distinguishable pairs.

If w distinguishes t from u , then sw distinguishes q from r .

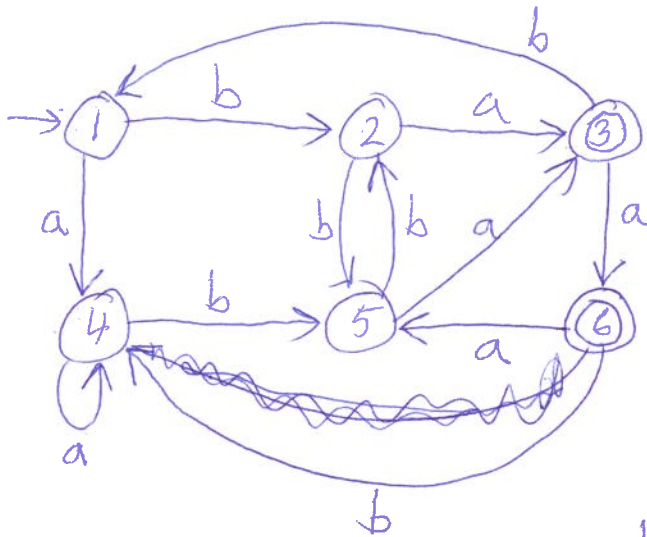
Merge 2,4,5 to get the minimum equivalent ⁽⁴⁾

DFA:



It will be clear and ~~unamb~~ unambiguous where to put arrows, start state, accepting states (no conflicts). If there is a conflict, that indicates an incorrect table ~~(X that should be~~ (omitting an X).

Ex:



2	X				
3	X	X			
4		X	X		
5	X		X	X	
6	X	X	X	X	X
	1	2	3	4	5

Merge

