# CSCE 574 ROBOTICS 

Path Planning

## Outline

- Path Planning
- Visibility Graph
- Potential Fields
- Bug Algorithms
- Skeletons/Voronoi Graphs
- C-Space


## Motion Planning

- The ability to go from $\mathbf{A}$ to $\mathbf{B}$
- Known map - Off-line planning
- Unknown Environment -Online planning
- Static/Dynamic Environment
- $\mathrm{q}_{\text {init }}$
- $q_{g o a l}$



## Path Planning

## World

Robot
Map

## Path Planning

## World

-Indoor/Outdoor
-2D/2.5D/3D
-Static/Dynamic
-Known/Unknown
Robot
-Abstract (web)
Map

## Path Planning

## World

## Robot

-Mobile
$>$ Indoor/Outdoor
$>$ Walking/Flying/Swimming
-Manipulator
-Humanoid
-Abstract

## Path Planning

## World



Map
-Topological

- Metric
-Feature Based
-1D,2D,2.5D,3D


## Path Planning

## World

-Indoor/Outdoor
-2D/2.5D/3D
-Static/Dynamic
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Robot
-Abstract (web)

## Map

-Topological

- Metric
-Feature Based
-1D,2D,2.5D,3D
-Manipulator
- Humanoid
-Abstract


## Path Planning: Assumptions

- Known Map
- Roadmaps (Graph representations)
- Polygonal Representation



## Visibility Graph

## - Connect Initial and goal locations with all the visible vertices



## Visibility Graph

- Connect initial and goal locations with all the visible vertices
- Connect each obstacle vertex to every visible obstacle vertex


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## Visibility Graph

- Connect initial and goal locations with all the visible vertices
- Connect each obstacle vertex to every visible obstacle vertex
- Remove edges that intersect the interior of an obstacle

- $\mathrm{q}_{\text {init }}$
- $\mathrm{q}_{\text {goal }}$


## Visibility Graph

- Connect initial and goal locations with all the visible vertices
- Connect each obstacle vertex to every visible obstacle vertex
- Remove edges that intersect the interior of an obstacle
- Plan on the resulting graph


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## Visibility Graph

- An alternative path
- Alternative name: "Rubber band algorithm"

- $q_{\text {init }}$
- $\mathrm{q}_{\text {goal }}$


## Major Fault

- Point robot
- Path planning
like that guarantees to hit the obstacles



## Path Planning

## Potential Field methods

- compute a repulsive force away from obstacles



## Local techniques

## Potential Field methods

- compute a repulsive force away from obstacles
- compute an attractive force toward the goal



## Local techniques

## Potential Field methods

- compute a repulsive force away from obstacles
- compute an attractive force toward the goal
$\rightarrow$ let the sum of the forces control the robot


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## Local techniques

## Potential Field methods

- compute a repulsive force away from obstacles
- compute an attractive force toward the goal
$\rightarrow$ let the sum of the forces control the robot



## Sensor Based Calculations



## Major Problem?

## Local Minima!



脊

## Simulated Annealing

- Every so often add some random force


## Limited-knowledge path planning

- Path planning with limited knowledge
- Insect-inspired "bug" algorithms

Goal


- known direction to goal
- otherwise local sensing walls/obstacles encoders
-"reasonable" world

1. finitely many obstacles in any finite disc
2. a line will intersect an obstacle finitely many times

## Not truly modeling bugs...

Insects do use several cues for navigation:

visual landmarks
polarized light
chemical sensing

neither are the current bugsized robots
they're not ears...
Other animals use information from
magnetic fields
electric currents
temperature

bacteria


## Bug Strategy

Insect-inspired "bug" algorithms


- known direction to goal •
- otherwise only local sensing walls/obstacles encoders


## "Bug 0" algorithm

1) head toward goal
2) follow obstacles until you can head toward the goal again
3) continue

## Does It Work?

## Bug 1

## Insect-inspired "bug" algorithms

- known direction to goal
- otherwise only local sensing walls/obstacles encoders



## "Bug 1" algorithm

1) head toward goal


## Bug 1

## Insect-inspired "bug" algorithms

- known direction to goal
- otherwise only local sensing walls/obstacles encoders



## "Bug 1" algorithm

1) head toward goal

2 ) if an obstacle is encountered, circumnavigate it and remember how close you get to the goal


## Bug 1

## Insect-inspired "bug" algorithms

- known direction to goal
- otherwise only local sensing walls/obstacles encoders



## "Bug 1" algorithm

1) head toward goal

2 ) if an obstacle is encountered, circumnavigate it and remember how close you get to the goal
3) return to that closest point (by wall-following) and continue

## Bug 1 analysis

Distance Traveled


What are bounds on the path length that the robot takes?

Available Information:
$\mathrm{D}=$ straight-line distance from start to goal
$\mathrm{P}_{\mathrm{i}}=$ perimeter of the $i$ th obstacle

Lower and upper bounds?
Lower bound:
Upper bound:

## Bug 1 analysis

Distance Traveled


What are bounds on the path length that the robot takes?

Available Information:
D = straight-line distance from start to goal
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Lower and upper bounds?
Lower bound: D
Upper bound:

## Bug 1 analysis

## Distance Traveled



What are bounds on the path
length that the robot takes?
Available Information:
$\mathrm{D}=$ straight-line distance from start to goal
$\mathrm{P}_{\mathrm{i}}=$ perimeter of the $i$ th obstacle

Lower and upper bounds?

> Lower bound: D

Upper bound: $\quad \mathrm{D}+1.5 \Sigma \mathrm{P}_{\mathrm{i}}$

How good a bound? How good an algorithm?

## A better bug?

Call the line from the starting point to the goal the s-line


## A better bug?

Call the line from the starting
point to the goal the s-line

"Bug 2" algorithm

1) head toward goal on the $s$-line

## A better bug?

Call the line from the starting
point to the goal the s-line


## "Bug 2" algorithm

1) head toward goal on the $s$-line
2) if an obstacle is in the way, follow it until encountering the sline again.

## A better bug?



## "Bug 2" algorithm

1) head toward goal on the $s$-line
2) if an obstacle is in the way, follow it until encountering the sline again.
3) Leave the obstacle and continue toward the goal

## A better bug?



## "Bug 2" algorithm

1) head toward goal on the $s$-line
2) if an obstacle is in the way, follow it until encountering the sline again closer to the goal.
3) Leave the obstacle and continue toward the goal

## Bug 2 analysis

## Distance Traveled



Goal

What are bounds on the path length that the robot takes?

Available Information:
D = straight-line distance from start to goal
$\mathrm{P}_{\mathrm{i}}=$ perimeter of the $i$ th obstacle

Lower and upper bounds?
Lower bound:
Upper bound:

## Bug 2 analysis

## Distance Traveled



Goal

What are bounds on the path
length that the robot takes?
Available Information:
$\mathrm{D}=$ straight-line distance from start to goal
$\mathrm{P}_{\mathrm{i}}=$ perimeter of the $i$ th obstacle
$\mathrm{N}_{\mathrm{i}}=$ number of s-line intersections with the $i$ th obstacle

Lower and upper bounds?
Lower bound:
Upper bound:

## Bug 2 analysis

## Distance Traveled



Goal

What are bounds on the path
length that the robot takes?
Available Information:
$\mathrm{D}=$ straight-line distance from start to goal
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Lower and upper bounds?
Lower bound: D
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## Bug 2 analysis

## Distance Traveled



Goal

What are bounds on the path
length that the robot takes?
Available Information:
D = straight-line distance from start to goal
$\mathrm{P}_{\mathrm{i}}=$ perimeter of the $i$ th obstacle
$\mathrm{N}_{\mathrm{i}}=$ number of s-line intersections with the $i$ th obstacle

Lower and upper bounds?
Lower bound: D
Upper bound:

$$
\mathrm{D}+0.5 \sum_{\mathrm{i}} \mathrm{~N}_{\mathrm{i}} \mathrm{P}_{\mathrm{i}}
$$

## head-to-head comparison

or thorax-to-thorax, perhaps

What are worlds in which Bug 2 does better than Bug 1 (and vice versa)?

## head-to-head comparison

What are worlds in which Bug 2 does better than Bug 1 (and vice versa)?

Bug 2 beats Bug 1



## Bug Mapping



## Other bug-like algorithms

The Pledge maze-solving algorithm


1. Go to a wall
2. Keep the wall on your right
3. Continue until out of the maze

## Other bug-like algorithms

The Pledge maze-solving algorithm


1) Go to a wall
2) Keep the wall on your right
3) Continue until out of the maze
int a[1817];main(z,p,q,r)\{for(p=80;q+p-80;p=2*a[p]) for(z=9;z--;)q=3\&(r=time(0)+r*57)/7,q=q?q-1?q-2?1-p\%79?-1:0:p\%79-77?1:0:p<1659?79:0:p>158?-79:0,q?!a[p+q*2]?a[p+=a[p+=q]=q]=q:0:0;for(;q++1817;)printf(q\%79?"\%c":"\%c\n"," \#"[!a[q-1]]);\}

IOCCC random maze generator


## Tangent Bug

- Limited Range Sensor
- Tangent Bug relies on finding endpoints of finite, continues segments of the obstacles



## Tangent Bug



## Contact Sensor Tangent Bug



1. Robot moves toward goal until it hits obstacle 1 at H 1
2. Pretend there is an infinitely small sensor range and the direction which minimizes the heuristic is to the right
3. Keep following obstacle until robot can go toward obstacle again
4. Same situation with second obstacle
5. At third obstacle, the robot turned left until it could not increase heuristic
6. D_followed is distance between M3 and goal, d_reach is distance between robot and goal because sensing distance is zero

## Limited Sensor Range Tangent-Bug



## Infinite Sensor Range Tangent Bug



## Known Map

Brushfire Transform



## The Wavefront Planner: Setup

| 7 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| 2 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 2 |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0 | 11 | 12 | 13 | 14 | 15 |

## The Wavefront in Action (Part 1)

- Starting with the goal, set all adjacent cells with "0" to the current cell + 1
- 4-Point Connectivity or 8-Point Connectivity?
- Your Choice. We'll use 8-Point Connectivity in our example

| 7 | $\mathbf{0}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 6 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 0 | 0 | 0 | 0 |
|  | 0 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 0 | 0 | 0 |
| 2 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 3 | 3 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 3 | $\mathbf{2}$ |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |

## The Wavefront in Action (Part 2)

- Now repeat with the modified cells
- This will be repeated until no 0 's are adjacent to cells with values $>=2$
- 0 's will only remain when regions are unreachable

| 7 | $\mathbf{0}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 6 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 0 | 0 | 0 | 0 |
| 3 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 0 | 0 | 0 | 0 |
| 2 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 4 | 4 | 4 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 4 | 3 | 3 |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 4 | 3 | $\mathbf{2}$ |
|  | 0 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |

## The Wavefront in Action (Part 3)

- Repeat

| 7 | $\mathbf{0}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 6 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 0 | 0 | 0 | 0 |
| 3 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 5 | 5 | 5 | 5 |
|  | $\mathbf{n}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 5 | 4 | 4 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 5 | 4 | 3 | 3 |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 5 | 4 | 3 | $\mathbf{2}$ |
|  | 0 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |
| 10 | 15 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

## The Wavefront in Action (Part 3)

- Repeat

| 7 | $\mathbf{0}$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 6 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 6 | 6 | 6 | 6 |
| 3 | 0 | 0 | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | 5 | 5 | 5 | 5 |
| 2 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 6 | 5 | 4 | 4 | 4 |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 6 | 5 | 4 | 3 | 3 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 6 | 5 | 4 | 3 | $\mathbf{2}$ |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |

## The Wavefront in Action (Part 3)

- Until Done
- 0's would only remain in the unreachable areas

| 7 | 18 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 9 | 9 | 9 | 9 | 9 | 9 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6 | 17 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 8 | 8 | 8 | 8 | 8 |
| 5 | 17 | 16 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 7 | 7 | 7 | 7 |
| 4 | 17 | 16 | 15 | 15 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 6 | 6 | 6 | 6 |
| 3 | 17 | 16 | 15 | 14 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 5 | 5 | 5 | 5 |
| 2 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 4 | 4 |
| 1 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 3 |
| 0 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 78 | 8 |  |  |  | 12 | 13 | 14 | 15 |

## The Wavefront in Action

- To find the shortest path, according to your metric, simply always move toward a cell with a lower number
- The numbers generated by the Wavefront planner are roughly proportional to their distance from the goal

| 7 | 18 |  |  |  |  |  |  |  |  |  | 9 | 9 | 9 | 9 | 9 | 9 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6 | 17 |  | 16 | 15 | 14 | 13 | 12 | 11 | 10 |  |  | 8 | 8 | 8 | 8 | 8 |
| 5 | 17 | 16 |  | 15 | 14 | 13 | 12 | 11 | 10 | 9 |  |  | 7 | 7 | 7 | 7 |
| 4 | 17 | 16 | 15 |  | 1 | 1 | 1 | 1 | 1 | 1 | 1 |  |  | 6 | 6 | 6 |
| 3 | 17 | 16 | 15 |  | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 5 |  | 5 | 5 |
| 2 | 17 | 16 | 15 | 14 |  | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 |  | 4 |
| 1 | 17 | 16 | 15 | 14 | 13 |  | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 |  |
| 0 | 17 | 16 | 15 | 14 | 13 | 12 |  |  |  |  |  |  |  |  |  |  |
|  | $\begin{array}{lll}0 & 1 & 2\end{array}$ |  |  | 3 | 4 | 5 | 6 |  |  | 1 |  | 1 | 2 | 13 |  | 15 |



## An alternative roadmap



## Voronoi diagrams



These line segments make up the Voronoi diagram for the four points shown here.

Solves the "Post Office Problem"

## Voronoi diagrams



These line segments make up the Voronoi diagram for the four points shown here.

Solves the "Post Office Problem"

or, perhaps, more important problems...

## Voronoi diagrams


"true" Voronoi diagram
(isolates a set of points)

generalized Voronoi diagram What is it?

## Voronoi diagrams

> Let $\mathrm{B}=$ the boundary of $\mathrm{C}_{\text {free }}$ Let $\mathbf{q}$ be a point in $\mathrm{C}_{\text {free }} \cdot(\bullet)$


## Voronoi diagrams

Let $\mathrm{B}=$ the boundary of $\mathrm{C}_{\text {free }}$.
Let $\mathbf{q}$ be a point in $\mathrm{C}_{\text {free }}$.


Define $\operatorname{clearance}(q)=\min \{|q-p|\}$, for all $p \in B$

## Voronoi diagrams

Let $\mathrm{B}=$ the boundary of $\mathrm{C}_{\text {free }}$. Let $\mathbf{q}$ be a point in $\mathrm{C}_{\text {free }}$.


Define clearance $(\mathrm{q})=\min \{|\mathrm{q}-\mathrm{p}|\}$, for all $\mathrm{p} \in \mathrm{B}$
Define $\operatorname{near}(\mathrm{q})=\{\mathrm{p} \in \mathrm{B}$ such that $|\mathrm{q}-\mathrm{p}|=\operatorname{clearance}(\mathrm{q})\}$

## Voronoi diagrams

## Evaluation

+ maximizes distance from obstacles
+ reduces to graph search
+ can be used in higher-dimensions
- nonoptimal
- real diagrams tend to be noisy

Let $\mathrm{B}=$ the boundary of $\mathrm{C}_{\text {free }}$.
Let $\mathbf{q}$ be a point in $\mathrm{C}_{\text {free }}$.


Define clearance $(\mathrm{q})=\min \{|\mathrm{q}-\mathrm{p}|\}$, for all $\mathrm{p} \in \mathrm{B}$
Define $\operatorname{near}(\mathrm{q})=\{\mathrm{p} \in \mathrm{B}$ such that $|\mathrm{q}-\mathrm{p}|=\operatorname{clearance}(\mathrm{q})\}$ q is in the Voronoi diagram of $\mathrm{C}_{\text {free }}$ if $\mid$ near $(\mathrm{q}) \mid>1$

## Generalized Voronoi Graph (GVG)



## Generalized Voronoi Graph (GVG)



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## Generalized Voronoi Graph (GVG)

-Access GVG


## Generalized Voronoi Graph (GVG)

- Access GVG
-Follow Edge



## Generalized Voronoi Graph (GVG)

- Access GVG •Home to the MeetPoint
-Follow Edge



## Generalized Voronoi Graph (GVG)

- Access GVG •Home to the MeetPoint
-Follow Edge •Select Edge



## GVG construction using sonar



- Nomadic Scout
- Sonar (GVG navigation)
- Camera with omni-directional mirror (feature detection)
- Onboard 1.2 GHz processor


## GVG construction using sonar



## GVG construction using sonar



## Slammer in Action



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## Removing Edges



## Meetpoint Detection

- 3o uncertainty ellipse of explored meetpoints
- Meetpoint degree (branching factor)
- Distances to local obstacles
- Relative angle bearings
- Edge signature
- Edge length
- Edge Curvature
- Vertex signal


## Ear-based Exploration



## Uncertainty Reduction



After Loop-closure



## Simulation



## Simulated Environment



## Real Environment



## Work Presented at IROS 2014

Ear-based Exploration on Hybrid Metric/Topological Maps
Q. Zhang, D. Whitney, F. Shkurti, and I. Rekleitis
 School of Computer Science, McGill University


## Voronoi applications



A retraction of a 3d object $==$ "medial surface"

in 2d, it's call a medial a.8

## skeleton $\hookleftarrow$ shape


curve evolution

where wavefronts collide

centers of maximal disks
again reduces a 2 d (or higher) problem to a question about graphs...

## skeleton $\hookleftarrow$ shape


curve evolution

where wavefronts collide

centers of maximal disks again reduces a 2 d (or higher) problem to a question about graphs...


## Problems



The skeleton is sensitive to small changes in the object's boundary.

- graph isomorphism (and lotsceffeqthedbseeph questions) : NP-complete


## Roadmap problems

If an obstacle decides to roll away... (or wasn't there to begin with)


